

# Toward to next-generation database

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TYPICAL  
IPHONE USER



HOW SHE SEES  
HERSELF



HOW SHE'S SEEN BY  
BLACKBERRY USERS



HOW SHE'S SEEN BY  
ANDROID USERS



TYPICAL  
ANDROID USER



HOW HE SEES  
HIMSELF



HOW HE'S SEEN BY  
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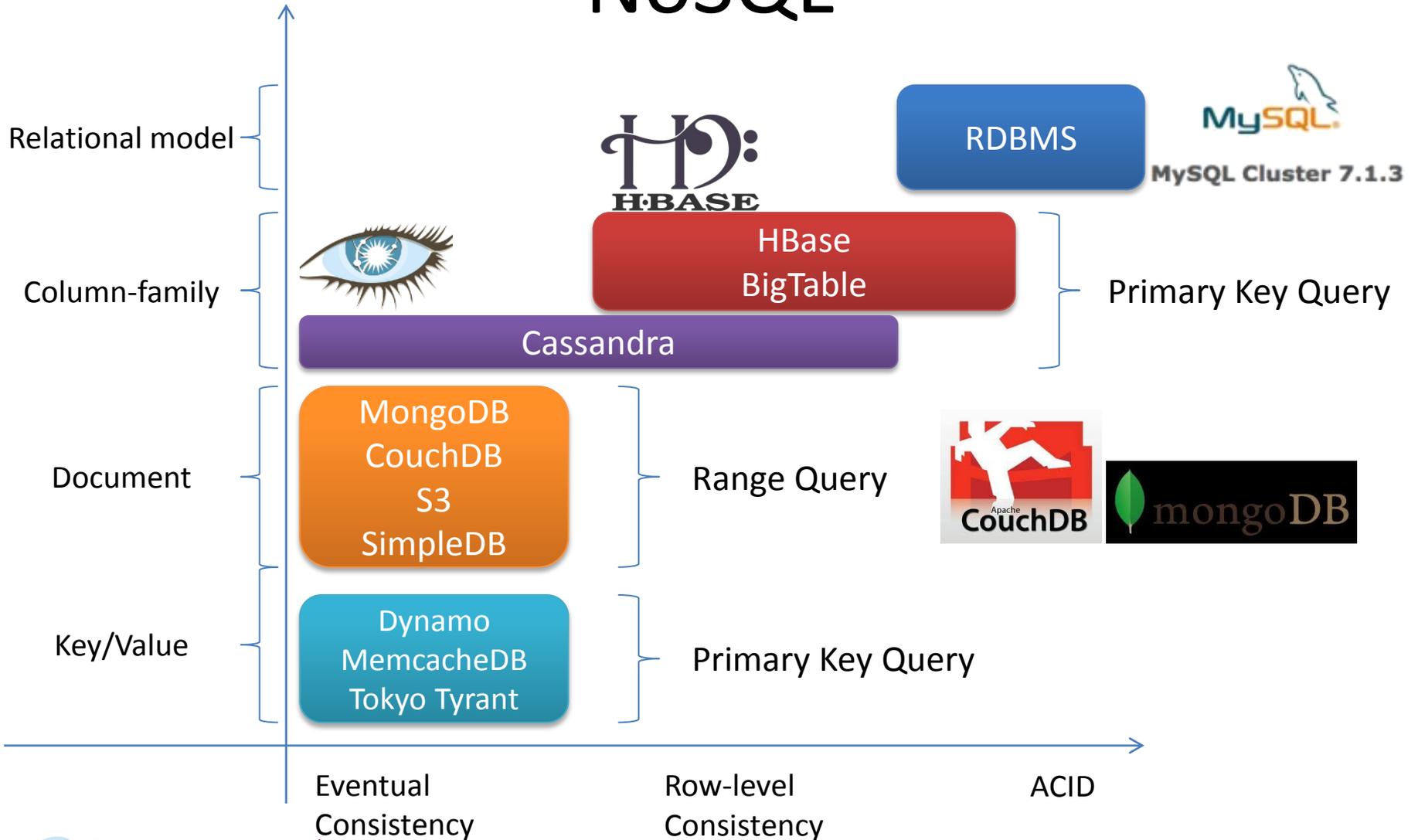
# Who am I

- A Master Student
  - HPDS Lab, Inst. of CCE, NCKU
  - 謝錫堃教授
- An Entrepreneur
  - PhiCloud.com
- Research Interest
  - Distributed File System
    - On-demand Data Co-allocation with User-Level Cache for Grids
      - Concurrency and Computation: Practice and Experience
  - MapReduce
    - Variable-Sized Map and Locality-Aware Reduce on Public-Resource Grids
      - GPC 2010
  - Distributed Transaction Processing
    - TSorter: A Conflict-Aware Transaction Processing System for Clouds
      - NCKU master thesis

# Outline

- NoSQL
- Existing Transactional NoSQL
- TSorter
- Hadoop/HBase Experience
- Future Work

# NoSQL



# Characteristics

- **Schema-Free**
- **Easy Replication Support**
- **Simple API**
- **Eventually Consistent**
- **BASE**
  - Basically Available, Soft State, Eventually Consistency
- **Huge Data Amount**

From <http://nosql-database.org/>



# Database Usage

- Analytical Processing
  - MapReduce Programming
  - Hive and Pig
- Transaction Processing
  - Lack support of ACID

# ACID

- Atomicity
  - Group multiple operations
- Consistency
  - ~~– Sequential/Eventual Consistency~~
  - Schema (However, NoSQL is schema-free)
- Isolation
  - Transaction serializability
- Durability
  - HDFS-200

# Transactional HBase

- Optimistic Concurrency Control (OCC)
  - No Lock
  - So, Good Luck
- 2-Phase Commit (2PC)
  - Distributed Transaction Processing

# Transactional HBase

T<sub>1</sub>:  
value = 1  
value = 2  
value = get(row id=123, column="value")  
value = value + 1  
put(key, value)  
commit()

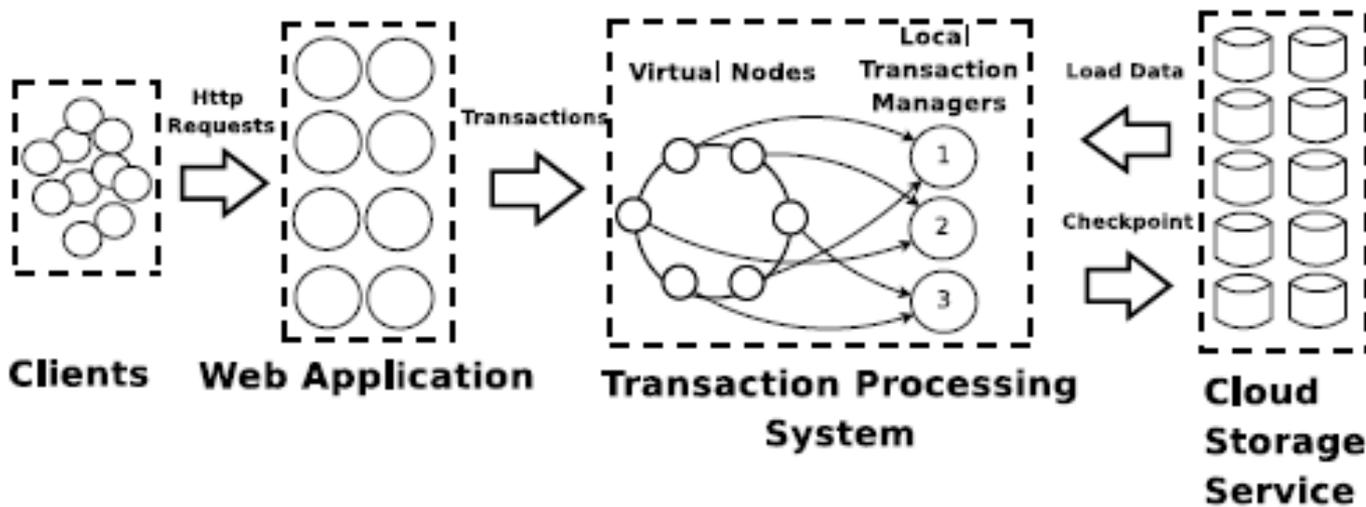
T<sub>2</sub>:  
value = 1  
value = 2  
value = get(row id=123, column="value")  
value = value + 1  
put(key, value)  
commit()



Row id	Values
...	...
123	{"value": 2}
...	...

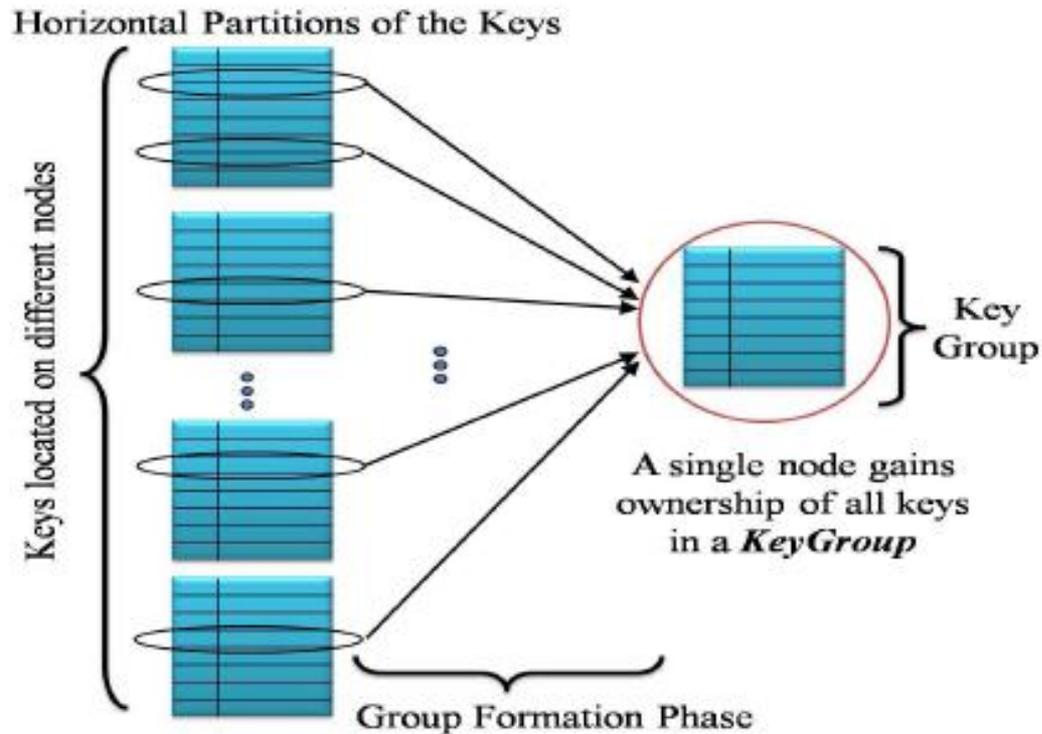
# CloudTPS

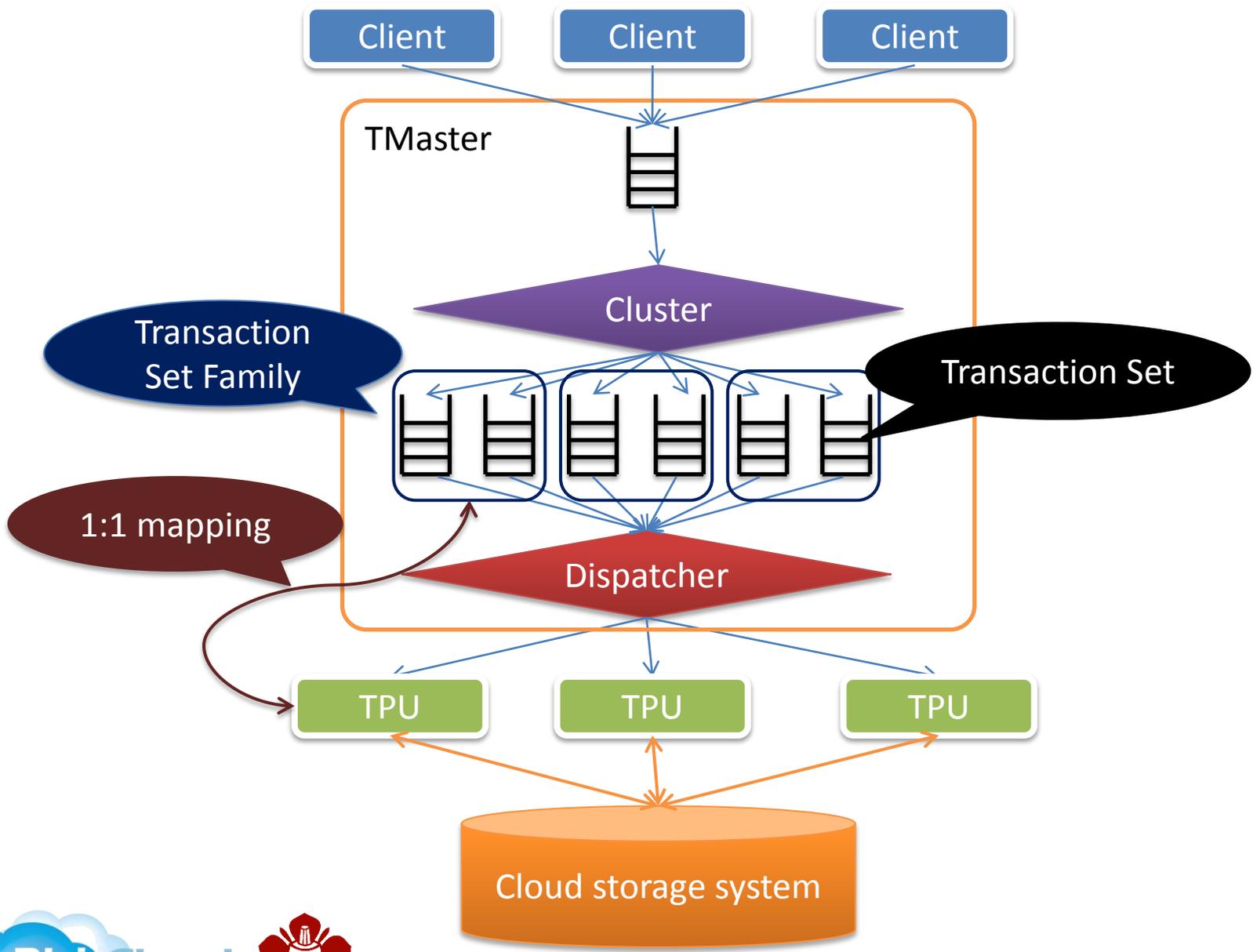
- Similar to Transactional HBase
- Timestamp-Based Concurrency Control
  - Sort sub-transactions by the global timestamp



# G-Store

- Key Grouping Protocol





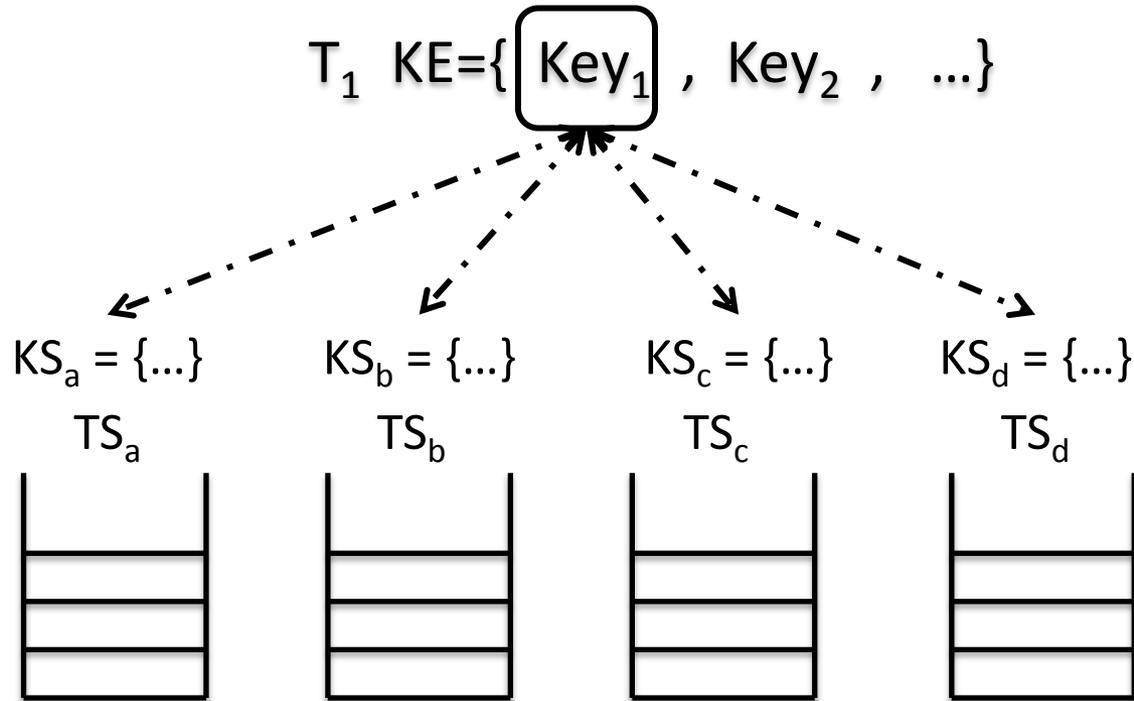
# Key Set

$T_a$ KE={Key <sub>1</sub> , Key <sub>4</sub> }
$T_b$ KE={Key <sub>1</sub> , Key <sub>3</sub> }
$T_c$ KE={Key <sub>1</sub> , Key <sub>2</sub> }

$TS_a$

$$KS_a = \{Key_1, Key_2, Key_3, Key_4\}$$

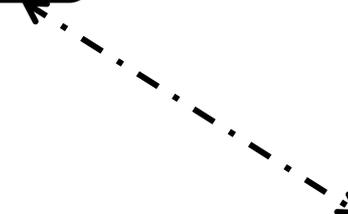
# Conflict Detection



*$O(\text{the number of keys} \times \text{the number of transaction sets})$*

# Conflict Detection

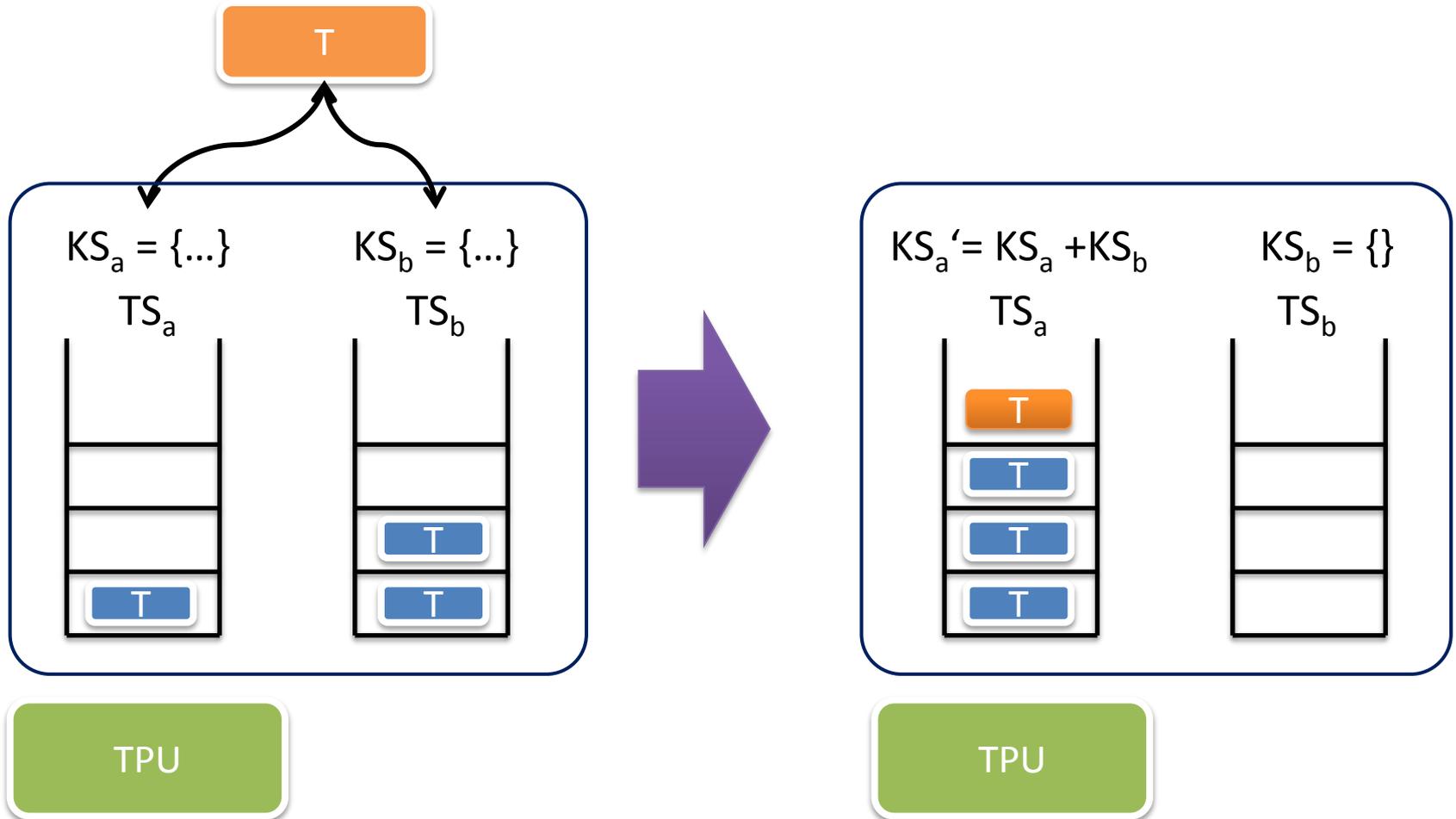
$T_1$  KE={ Key<sub>1</sub> , Key<sub>2</sub> , ... }



Key	Transaction set
Key <sub>1</sub>	TS <sub>a</sub>
Key <sub>2</sub>	TS <sub>b</sub>
Key <sub>3</sub>	TS <sub>c</sub>
Key <sub>4</sub>	TS <sub>d</sub>
...	...

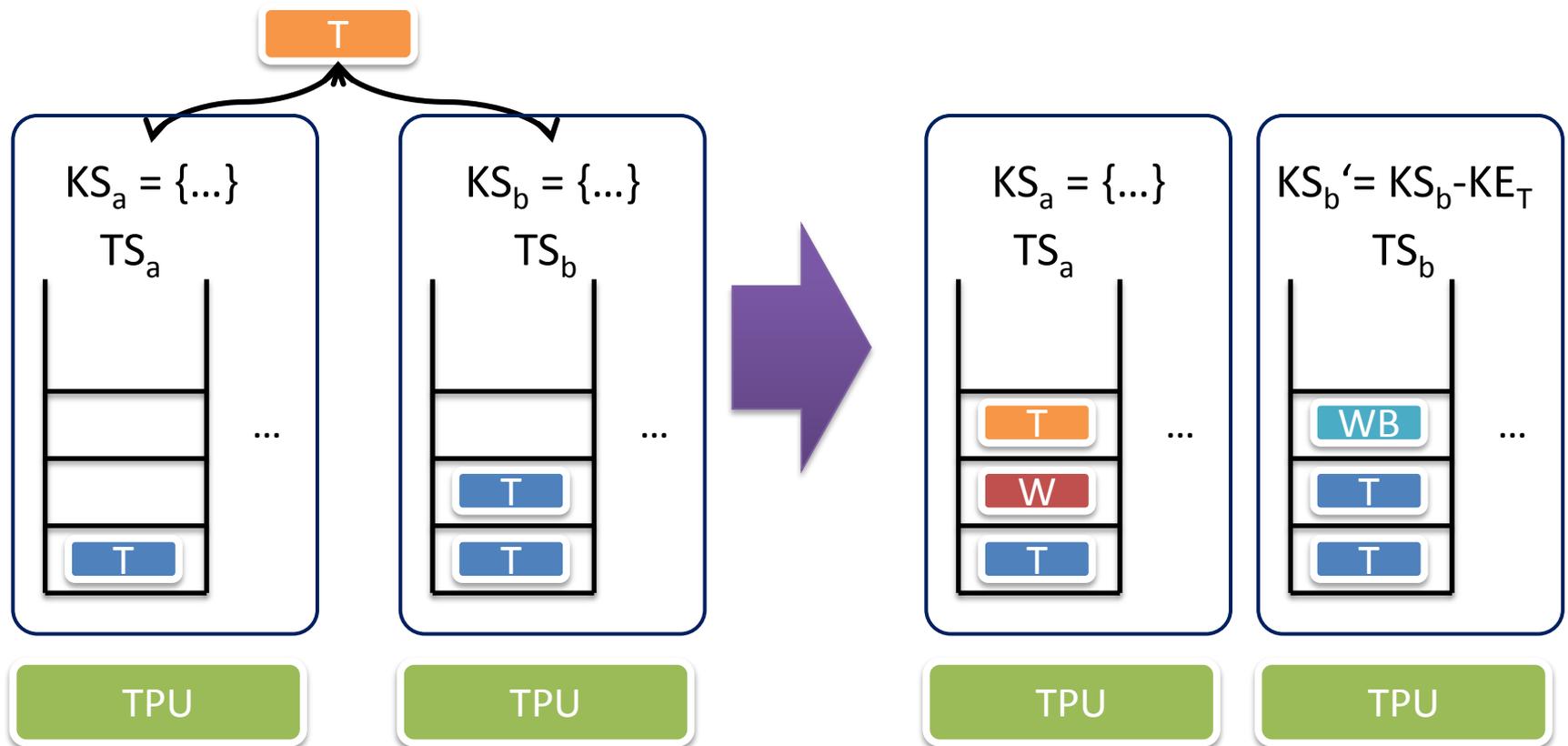
$O(\text{the number of keys})$

# Transaction Set Serialization

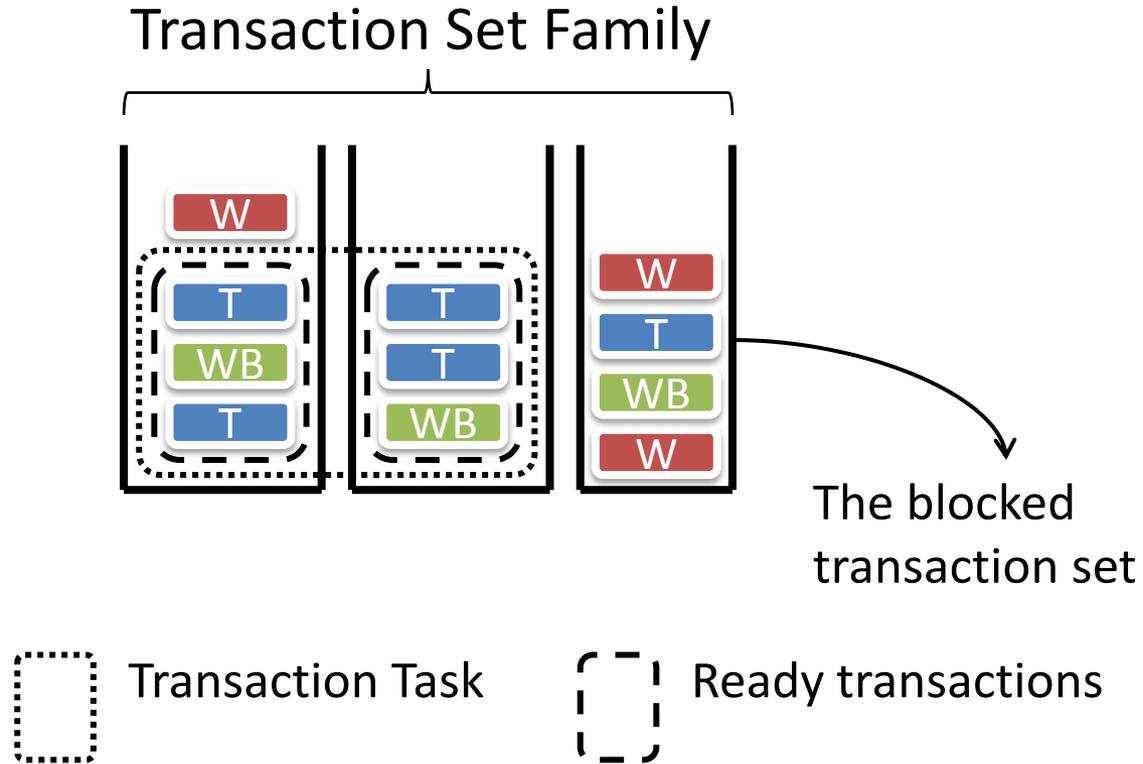


# Transaction Set Serialization

- Write-back transaction
- Wait transaction

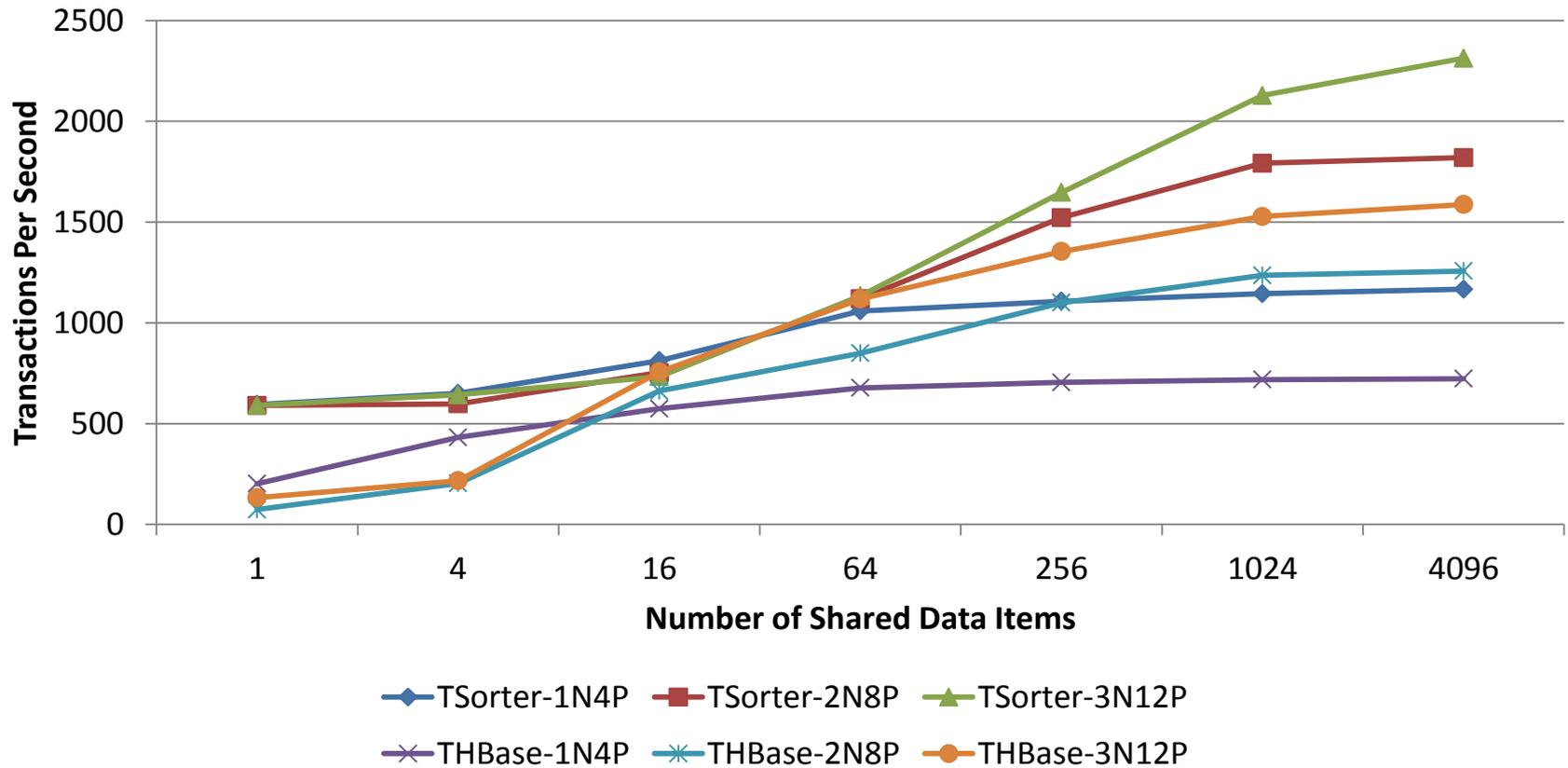


# Transaction Dispatcher



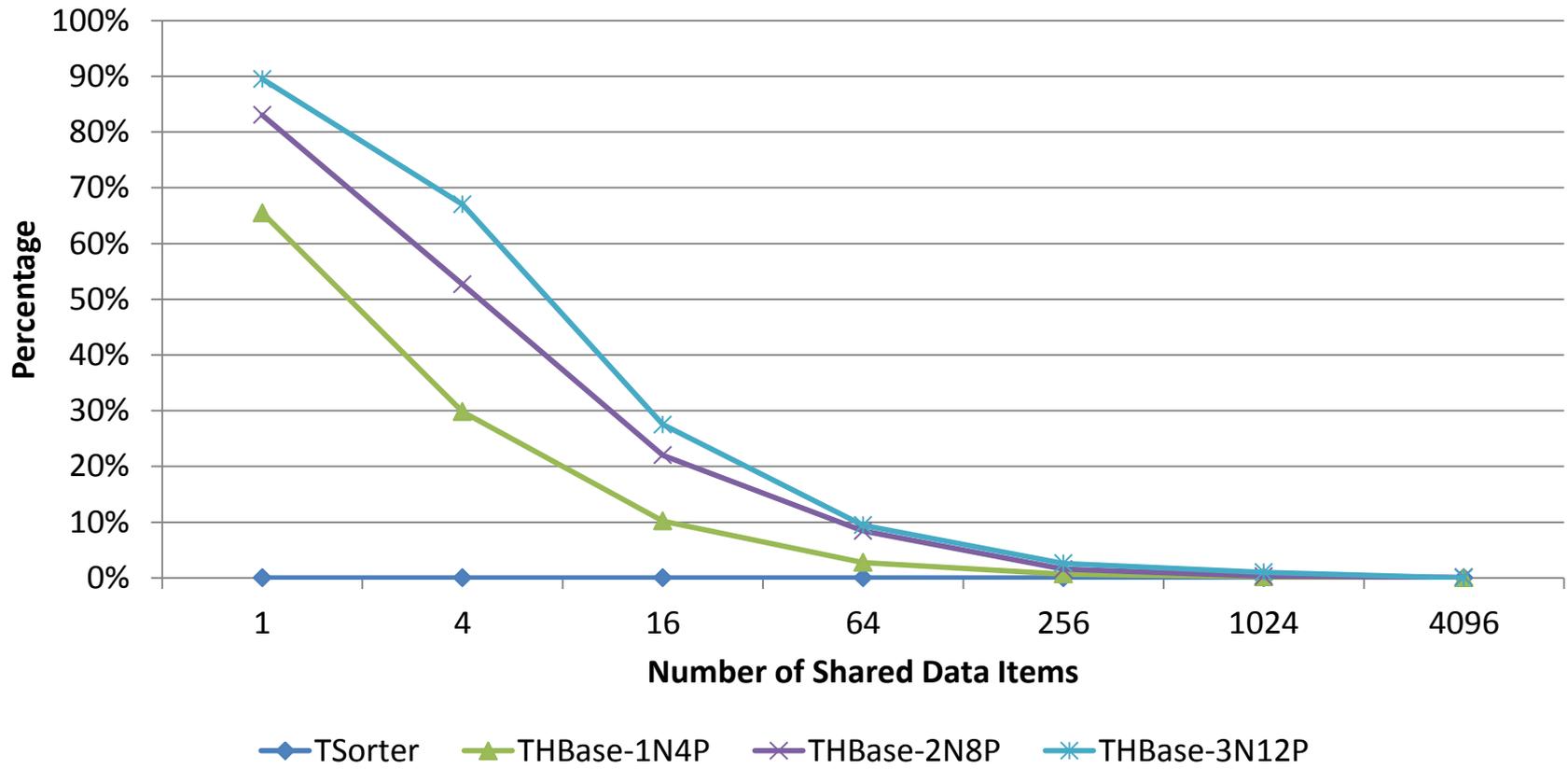
# Experiment

## Throughput of Conflict-Intensive Workload



# Experiment

## Abort Rate of Conflict-Intensive Wrokload



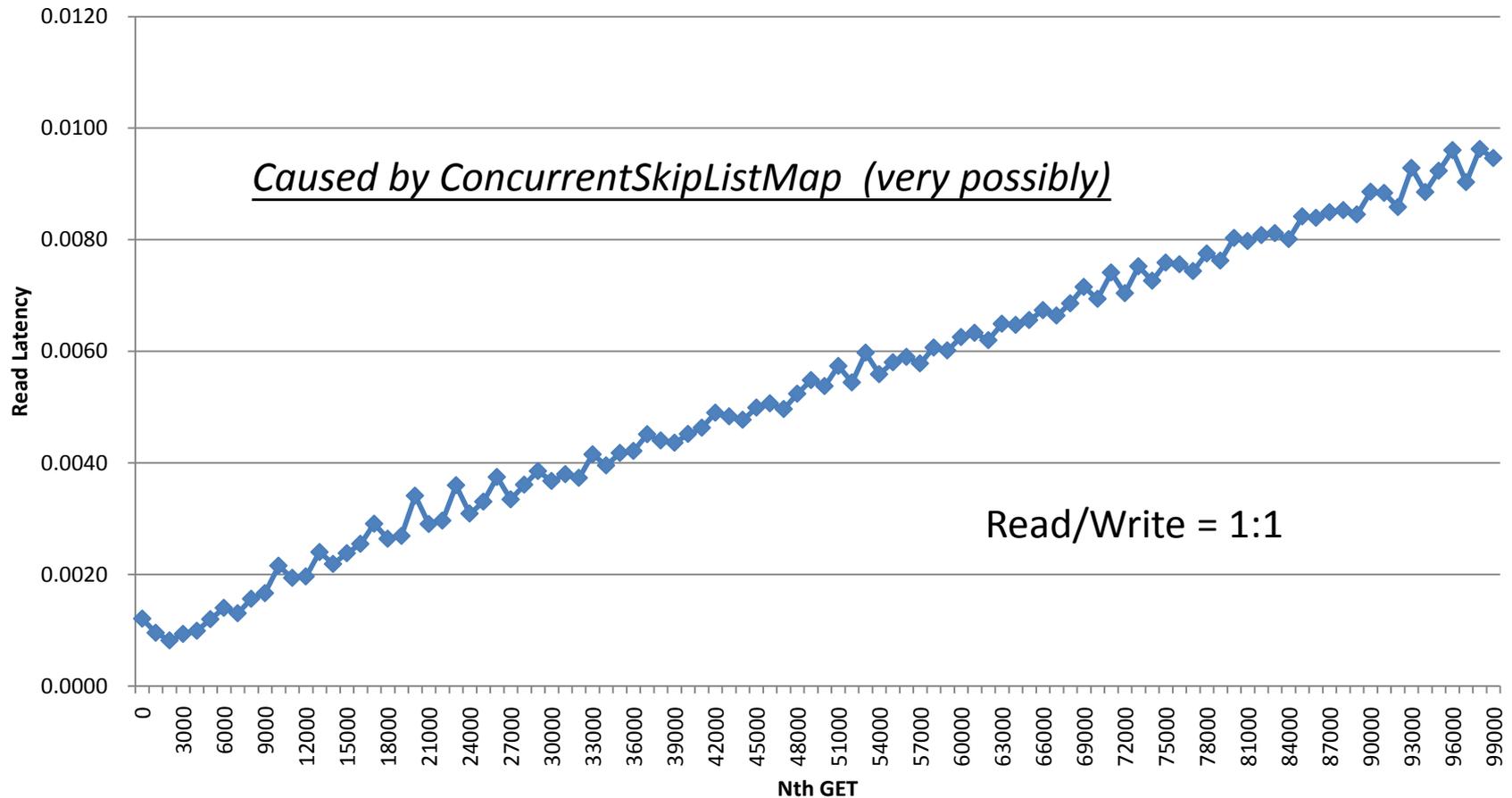
# Hadoop/HBase Experience

- Hadoop
  - Q1: HDFS as a Generic File System
- HBase
  - Q2: High Latency Read of HBase
  - Q3: HBase as a Web-Backend
  - Q4: HBase vs Cassandra

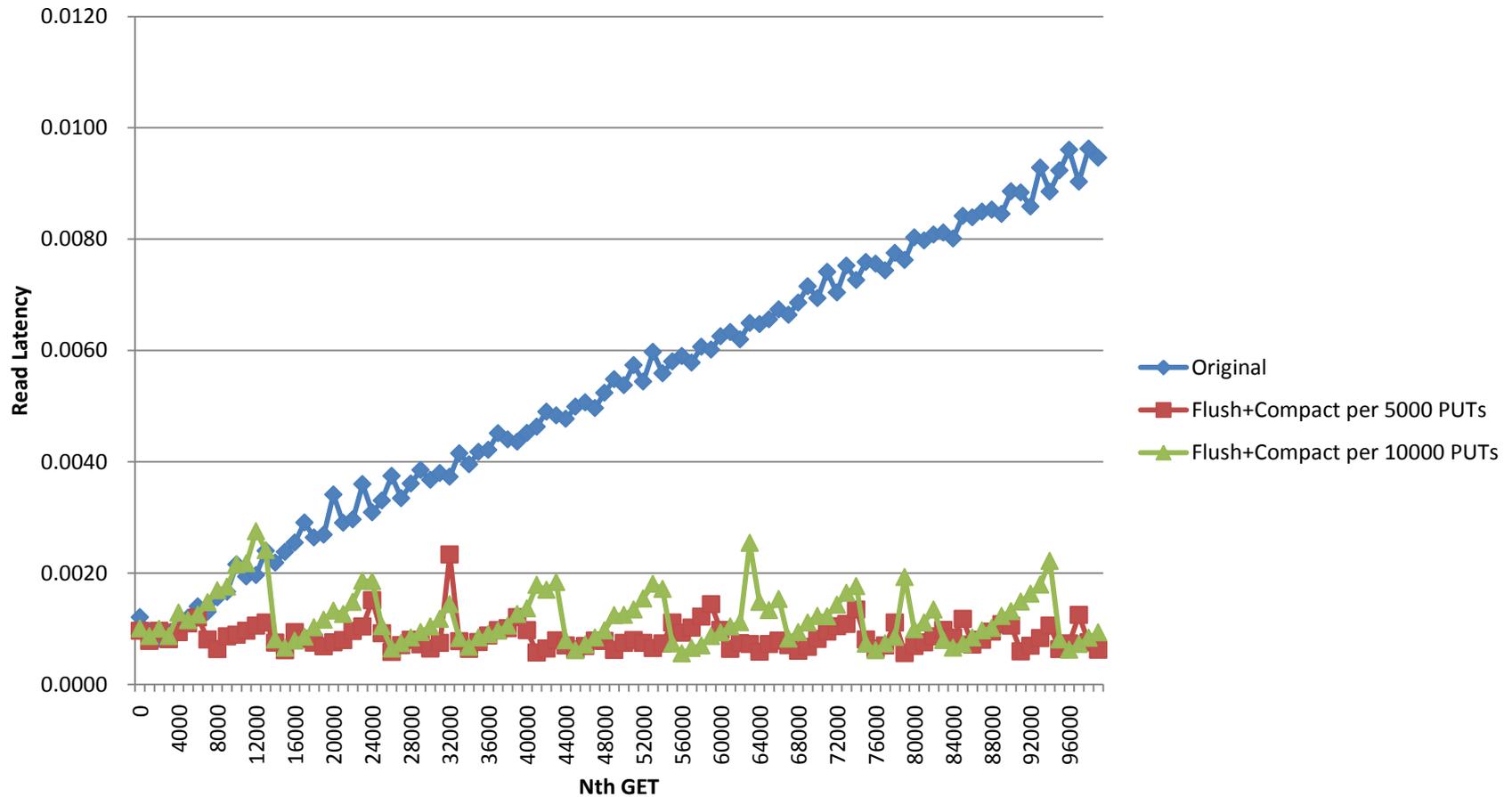
# HDFS as a Generic File System

- HDFS is not a POSIX-compliant FS
  - Write-Once-Read-Many
- You can Try GlusterFS
  - GlusterFS is one of the most sophisticated file system in terms of features and extensibility.
  - Features
    - Data Replicated/Striped
    - No metadata
    - Client-Side Implementation

# High Latency Read of HBase



# High Latency Read of HBase



# HBase as a Web-Backend

- ***wbbs.cc.ncku.edu.tw***
  - ~750,000 Rows
  - ~1.5 GB
    - ~50 MB Index
  - ~3000 requests per day
  - I use MySQL
- **力可科技**
  - ~ 1 Billion Rows
  - ~55 GB
    - ~24 GB Index
  - They use Cassandra
- Size
  - GB? TB? PB?
- Requests per Second
- Database Consolidation
  - Very Large Migration Cost

# HBase vs Cassandra

	HBase	Cassandra
Storage Layer	HDFS	No
Data Partition	Size-Based	Random
Data Cache	Block-Based (1GB=8192 Blocks)	Row-Based
Consistency	Row-Level Serialization	Eventual Consistency <ul style="list-style-type: none"><li>• Write-One-Read-All</li><li>• Write-All-Read-One</li><li>• <math>(W+R) &gt; (N/2)</math></li></ul>
Suitable for	<ul style="list-style-type: none"><li>• Analytical Processing<ul style="list-style-type: none"><li>• Search Engine</li></ul></li><li>• OLTP with the Strong Need of Serialization<ul style="list-style-type: none"><li>• FB new IM System</li></ul></li></ul>	OLTP <ul style="list-style-type: none"><li>• Mail</li><li>• Log</li><li>• Content Repo.</li></ul>

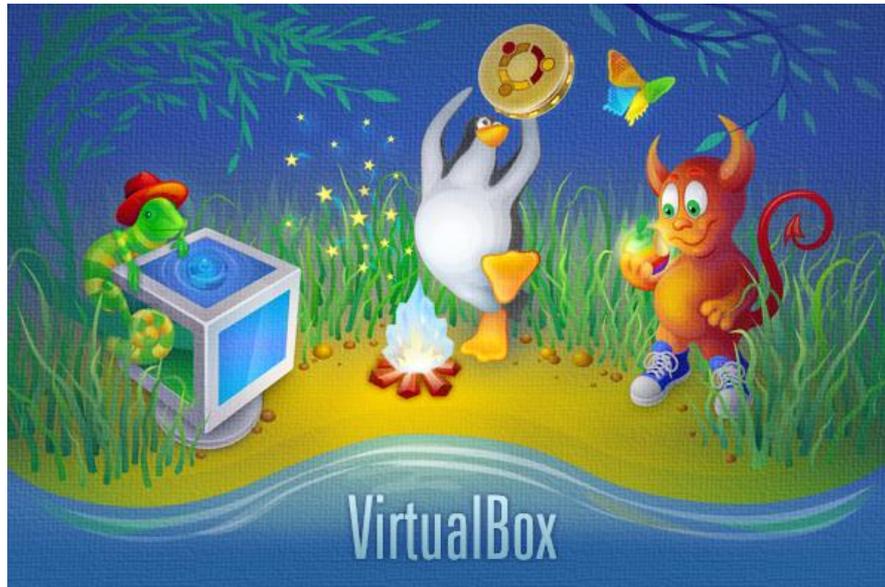
# Scalability ?

# ~~Big and Large~~

# From Small To Large

# PhiCloud

- **PhiCenter**
  - 2011 Q1
  - A System Like Eucalyptus, OpenNebula, OpenStack and XenServer
  - Easy-to-Use, Based-on VirtualBox-OSE and GPLed



# PhiCloud

- **PhiDesktop**
  - 2011 Q2
  - A Product Like VMWare View or XenDesktop
  - Address on University
    - 校園授權軟體
  - 50% Features and 20% Costs
    - V牌 USD150
    - C牌 USD100
  - Components
    - VirtualBox-OSE as Hypervisor, PhiCenter as VM Management
    - Xvfb as X11 Environment
    - XDamage Extension for Damage Detection
    - libjpeg-turbo for Image Compression

# PhiCloud

- **PhiData**
  - 2011 Q2
  - Public Cloud of Data Warehouse as a Service
  - Hive-Based or Develop a new MapReduce Engine
  - Welcome to Join PhiCloud

